

# Secrets of Elliptic Curves

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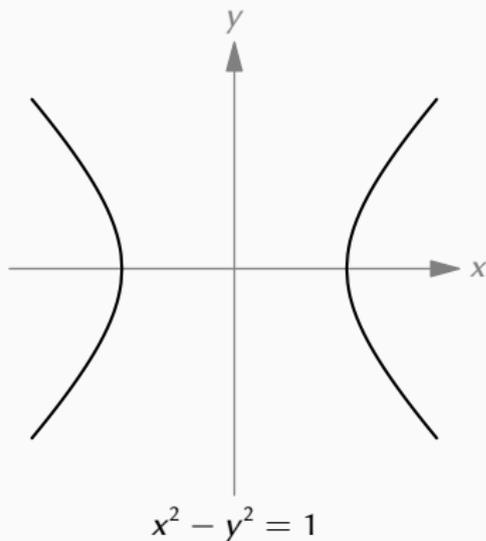
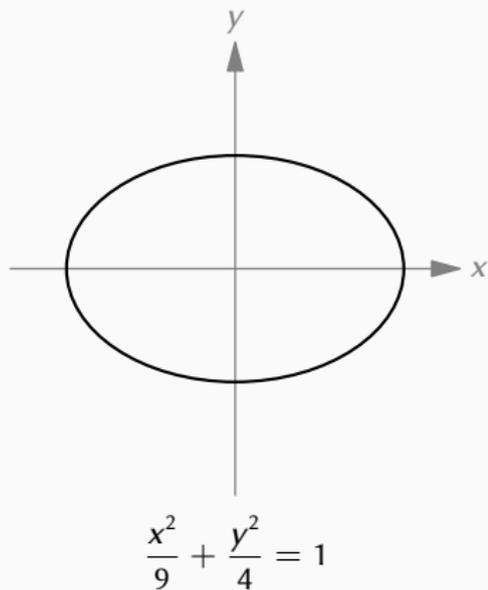
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October 24, 2025

# Conics

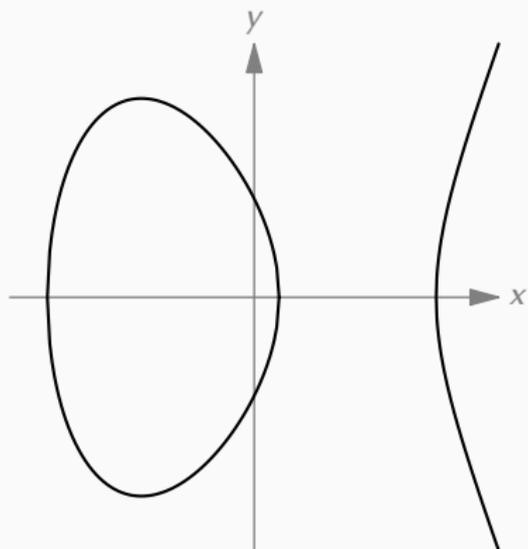
In high school, we learned about conics, which are curves described by degree 2 equations.



# Elliptic Curve

Today, we are going to talk about elliptic curves, which are curves of the form

$$y^2 = x^3 + ax + b.$$

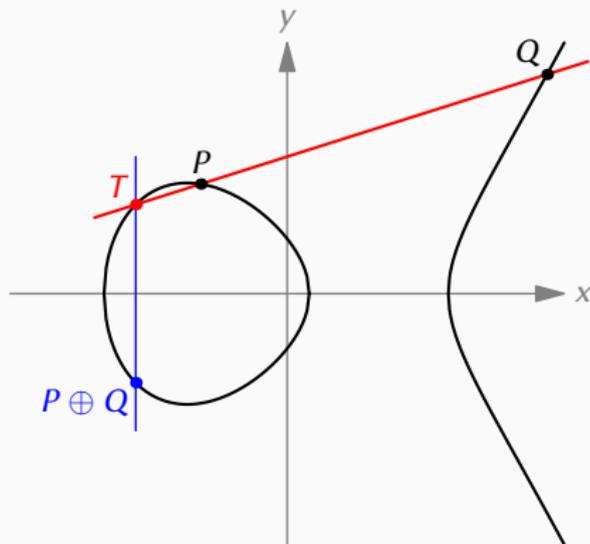


$$y^2 = x^3 - 4x + 1$$

# Group Operation

Given points  $P$  and  $Q$  on the elliptic curves, we **add** those two points by

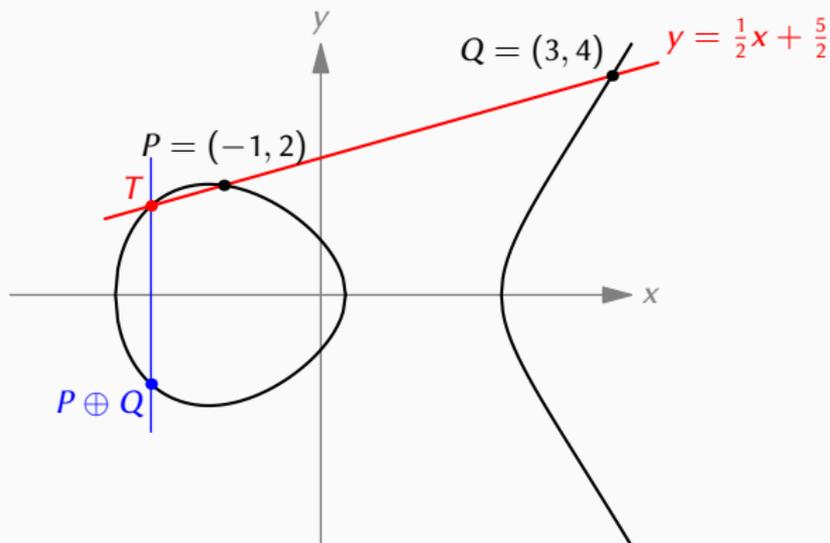
- Draw the line through point  $P$  and  $Q$ .
- This line intersects the elliptic curve at one other point. Let this point be  $T$ .
- Reflect  $T$  across the  $x$ -axis to get point  $P \oplus Q$ .



## Example

Let  $E$  be the curve  $y^2 = x^3 - 4x + 1$ ,  $P = (-1, 2)$ , and  $Q = (3, 4)$ .

- The line through  $P$  and  $Q$  has equation  $y = \frac{1}{2}x + \frac{5}{2}$ .
- This line intersect the elliptic curve again at  $T = (-\frac{7}{4}, \frac{13}{8})$ .
- Therefore,  $P \oplus Q = \boxed{(-\frac{7}{4}, -\frac{13}{8})}$ .



## Do order matter?

If we have three points  $P$ ,  $Q$ , and  $R$ , we can add them in two ways:

$$(P \oplus Q) \oplus R \quad \text{and} \quad P \oplus (Q \oplus R).$$

### Example

In the curve  $y^2 = x^3 - 4x + 1$ , if

$$P = (-1, 2), \quad Q = (3, 4), \quad R = (0, 1).$$

Then

$$\begin{aligned} P \oplus Q &= \left(-\frac{7}{4}, -\frac{13}{8}\right), & (P \oplus Q) \oplus R &= \left(\frac{92}{49}, \frac{113}{343}\right) \\ Q \oplus R &= \left(-\frac{2}{9}, \frac{37}{27}\right), & P \oplus (Q \oplus R) &= \left(\frac{92}{49}, \frac{113}{343}\right). \end{aligned}$$

## Theorem (Associativity of $\oplus$ )

*For any points  $P, Q, R$  on an elliptic curve, we have*

$$(P \oplus Q) \oplus R = P \oplus (Q \oplus R).$$

One can prove this by deriving the formula for  $P \oplus Q$ . Let's see how it goes.

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Let  $P = (x_P, y_P)$ ,  $Q = (x_Q, y_Q)$ ,  $R = (x_R, y_R)$ . Then the line through  $P$  and  $Q$  has equation

$$y - y_P = \frac{y_P - y_Q}{x_P - x_Q}(x - x_P),$$

which intersects the elliptic curve again at

$$\left( \left( \frac{y_P - y_Q}{x_P - x_Q} \right)^2 - x_P - x_Q, \left( \frac{y_P - y_Q}{x_P - x_Q} \right)^3 + \frac{y_P + y_Q}{2} - \frac{3}{2} \cdot \frac{y_P - y_Q}{x_P - x_Q} (x_P + x_Q) \right).$$

Negating the  $y$ -coordinate gives point  $P \oplus Q$ .

## Theorem (Associativity of $\oplus$ )

For any points  $P, Q, R$  on an elliptic curve, we have

$$(P \oplus Q) \oplus R = P \oplus (Q \oplus R).$$

Thus, we have that  $P \oplus Q$  is the point

$$\left( \left( \frac{y_P - y_Q}{x_P - x_Q} \right)^2 - x_P - x_Q, - \left( \frac{y_P - y_Q}{x_P - x_Q} \right)^3 - \frac{y_P + y_Q}{2} + \frac{3}{2} \cdot \frac{y_P - y_Q}{x_P - x_Q} (x_P + x_Q) \right).$$

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We now need to perform the same addition operation with  $R$ , getting a messy expression for the coordinates of  $(P \oplus Q) \oplus R$ .

Similarly, we can get another messy expression for  $P \oplus (Q \oplus R)$ .

These expressions have about 500 terms. However, they miraculously turn out to be equal!

## Why do we care?

The property that we just proved allows us to add points in any order. For example,

$$P \oplus P = 2P$$

$$2P \oplus 2P = 4P$$

$$4P \oplus 4P = 8P$$

$$\vdots$$

(where  $nP = \underbrace{P \oplus \cdots \oplus P}_{n \text{ } P\text{'s}}$ ). This allows us to compute  $nP$  for large  $n$  quickly.

On the other hand, the relation between  $P$  and, e.g.,  $1000P$  is very complicated. Given  $nP$  and  $P$ , no one knows how to efficiently determine  $n$ .

This hardness is a foundation of **elliptic curve cryptography**.

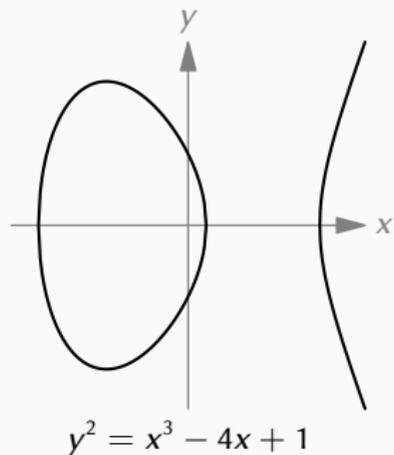
# Is this the end?

Our proof of  $(P \oplus Q) \oplus R = P \oplus (Q \oplus R)$  raises some burning questions.

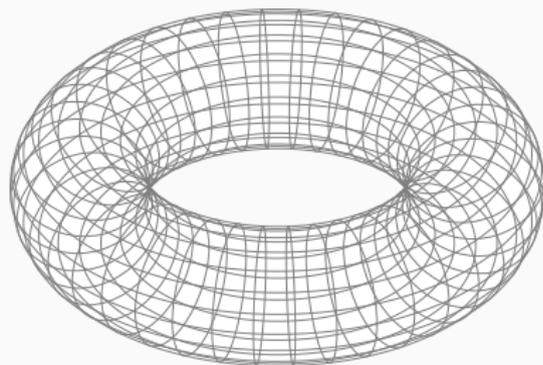
- Is there any way to avoid these massive computations?
- Why should we expect this to be true?

Mathematicians have answer to these questions through two different perspectives.

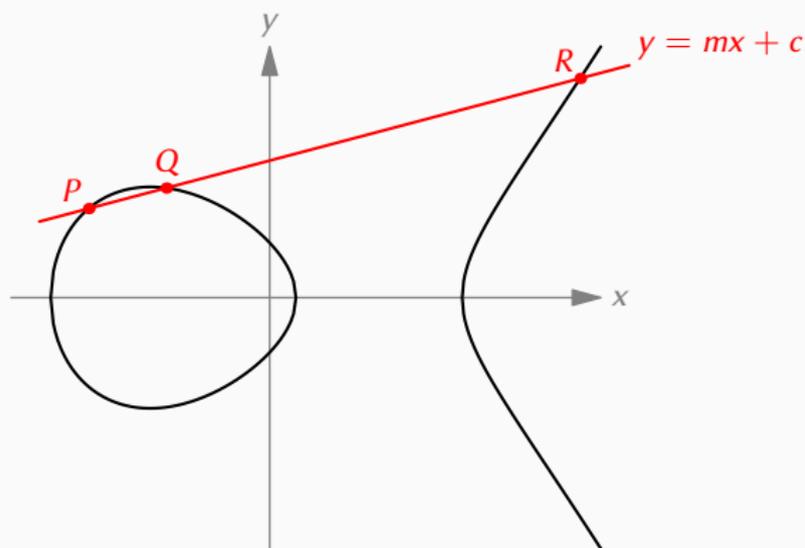
# Perspective 1: Complex Analysis



$\mathbb{R}^2$



## Perspective 2: Divisors

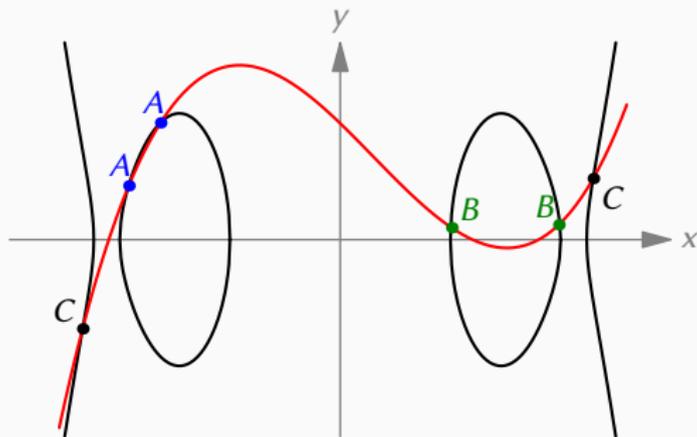


$$\operatorname{div}(y - mx - c) = (P) + (Q) + (R).$$

## More Complicated Curves

Either of these two perspectives allow us to **see** how to generalize the addition law to more complicated curves.

For example, the curve  $y^2 = x^6 + ax^5 + bx^4 + \cdots + ex + f$  have an addition law on **pairs** of points.



$$A + B + C = 0 \iff \text{there is a cubic passing through } A, B, C$$

Any curve have an addition law of  $g$ -tuples of points, where  $g$  is the **genus**, which measures complexity of a curve.

# Why do we study advanced mathematics?

Advanced mathematics can **abstract away** complicated equations, allowing us to see things (e.g., elliptic curves addition law) in a simpler ways.

With new understanding, we can reveal hidden structure that we haven't seen before (e.g., higher-genus curves).

**Thank You!**