

On t -intersecting Families of Permutations

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Intersecting Families of Permutations

Consider a collection \mathcal{F} of permutations of $[n] := \{1, 2, \dots, n\}$.

We say that \mathcal{F} is **t -intersecting** if any two permutations in \mathcal{F} agree at $\geq t$ inputs, i.e., for any $\sigma_1, \sigma_2 \in \mathcal{F}$, there exist $\geq t$ indices $i \in [n]$ such that $\sigma_1(i) = \sigma_2(i)$.

Question

For each $t < n$, what is the maximum size of t -intersecting family of permutations of $[n]$?

Construction 1

The family

$$\mathcal{A}_0 = \{\sigma : \sigma(i) = i \text{ for all } i \in \{1, 2, \dots, t\}\}$$

is clearly t -intersecting and has size $(n - t)!$.

For example, if $n = 7$ and $t = 4$, the elements in \mathcal{A}_0 are (in one-line notation)

$$\begin{array}{l} (1, 2, 3, 4, \mathbf{5}, \mathbf{6}, \mathbf{7}) \quad (1, 2, 3, 4, \mathbf{5}, \mathbf{7}, \mathbf{6}) \\ (1, 2, 3, 4, \mathbf{6}, \mathbf{7}, \mathbf{5}) \quad (1, 2, 3, 4, \mathbf{6}, \mathbf{5}, \mathbf{7}) \quad . \\ (1, 2, 3, 4, \mathbf{7}, \mathbf{5}, \mathbf{6}) \quad (1, 2, 3, 4, \mathbf{7}, \mathbf{6}, \mathbf{5}) \end{array}$$

This family turns out to be optimal when t is not too close to n .

Construction 2

The family

$$\mathcal{A}_k = \{ \sigma : \sigma(i) = i \text{ for } \geq t+k \text{ indices } i \text{ in } \{1, 2, \dots, t+2k\} \}$$

is clearly t -intersecting.

For example, if $n = 7$, $t = 4$, and $k = 1$, the elements in \mathcal{A}_1 are (in one-line notation)

$$\begin{aligned} & (1, 2, 3, 4, 5, 6, 7) \\ & (7, 2, 3, 4, 5, 6, 1) \quad (1, 7, 3, 4, 5, 6, 2) \\ & (1, 2, 7, 4, 5, 6, 3) \quad (1, 2, 3, 7, 5, 6, 4) \\ & (1, 2, 3, 4, 7, 6, 5) \quad (1, 2, 3, 4, 5, 7, 6) \end{aligned} ,$$

so $|\mathcal{A}_1| = 7$, larger than $|\mathcal{A}_0| = 6$.

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Conjecture (Deza and Frankl; Cameron)

Any t -intersecting family of permutations has size at most $\max_k |\mathcal{A}_k|$.

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This conjecture holds for the following ranges of n and t :

Range	Year	Authors	Method
$t \geq n - \omega(1)$	1977	Deza, Frankl	Elementary
$t = 1$	2003	Cameron, Ku	Elementary
$t \leq O(\log \log n)$	2011	Ellis, Friedgut, Pilpel	Rep. theory
$t \leq O\left(\frac{\log n}{\log \log n}\right)$	2022	Ellis, Lifshitz	Fourier + Rep. theory
$t \leq O\left(\frac{n}{\log^2 n}\right)$	2024	Kupavskii, Zakharov	Spread Approx.
$t \leq O(n)$	2024	Keller, Lifshitz, Minzer, Sheinfeld	Fourier
$t \leq (1 - o(1))n$	2024+	Kupavskii	Spread Approx.

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Theorem (S., 2025+)

This conjecture holds for $t \leq n - O(n^{3/4+\varepsilon})$ for any $\varepsilon > 0$.

Partial Permutations

In the proof, it's more helpful to view a permutation as a subset of $[n]^2$.
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A **partial permutation** is a subset of $[n]^2$ that is contained in some permutation. An example of a partial permutation is

$$\{(1, 4), (3, 2), (6, 7)\}.$$

The **length** of a partial permutation is the size of the set (so the example above has length 3).

Kupavskii's Proof for $t \leq (1 - o(1))n$

Let \mathcal{F} be a t -intersecting family of permutations.

1. (Spread Approximation) Construct a set of partial permutations \mathcal{S} such that
 - each element of \mathcal{S} has length $\leq t + 0.001(n - t)$.
 - \mathcal{S} is t -intersecting: any two elements have intersection of size $\geq t$.
 - $\geq 99.9\%$ of elements in \mathcal{F} contains at least one element of \mathcal{S} .

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2. **(Peeling)** Keep removing entries from a partial permutation in \mathcal{S} while as long as \mathcal{S} is still t -intersecting.

In the resulting set, let \mathcal{W}_k be the set of partial permutations of length $t + k$. Choose k such that $|\mathcal{W}_k| \geq t^{-0.5} \binom{t}{k}$ (i.e., large).

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3. **(Counting)** There exists $A, B \in \mathcal{W}_k$ such that $|A \cap B| = t$. Show that 99% of elements in \mathcal{F} contains an element in $\binom{A \cup B}{t+k}$.

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4. **(Finishing)** Suppose that there exists an element not containing a set in $\binom{A \cup B}{t+k}$. Obtain a final contradiction by counting.

Our Improvement

Step 2, 3, and 4 can be easily modified to work if $t \leq n - O(n^{3/4+\varepsilon})$.

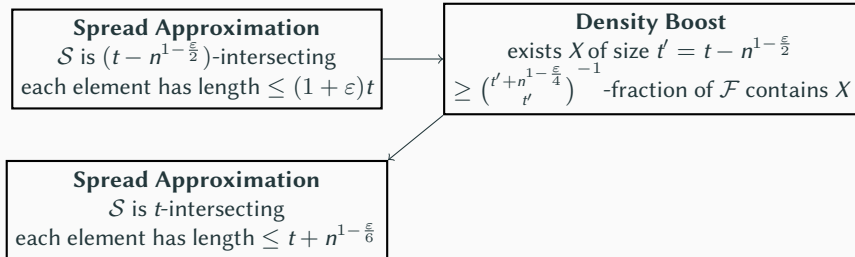
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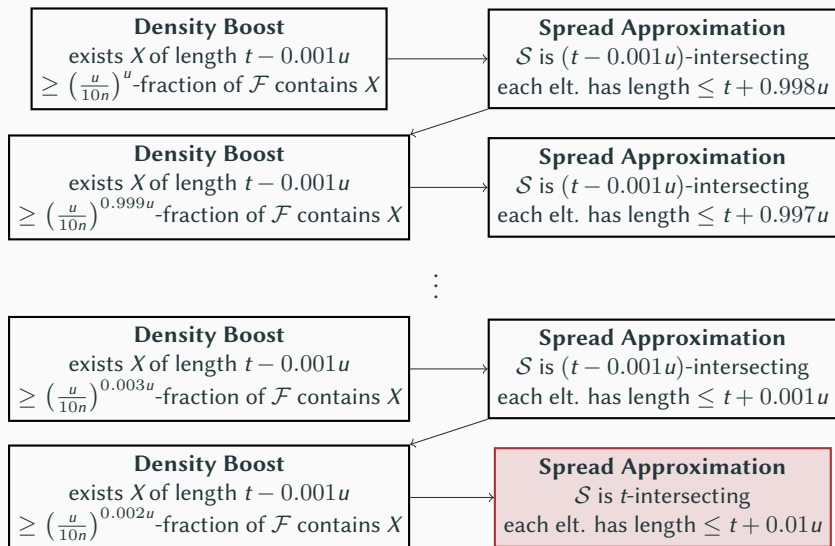
Kupavskii's proof works in the following way.



The key tool of the spread approximation step is the **sunflower lemma**, which is based on Alweiss-Lovett-Wu-Zhang's result and sharpened by Tao. This limits the bound that we can get.

Our Proof: Iterated Spread Approximation

Let $u = n - t$ and assume $u = \Omega(n^{0.6+\varepsilon})$.



Thank you!

Towards the Complete Intersection Theorem for Permutations

The new Step 1 works for $t \leq n - O(n^{1/2+\epsilon})$.

We now need to look at Step 3 (the counting step).

Original Proof

- Prove that there exists $A, B \in \mathcal{W}_k$ (with $t + k$ elements) such that $|A \cap B| = t$.
- Count the number of sets in \mathcal{W}_k that intersect both A and B with size $\geq t$.
- Result: most elements in \mathcal{W}_k are contained in $A \cup B$ (size $t + 2k$).

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Potential Improvement

- Prove that there exists $A_1, \dots, A_r \in \mathcal{W}_k$ (with $t + k$ elements) such that **[insert a condition]**
- Count the number of sets in \mathcal{W}_k that intersect each of A_1, \dots, A_r with size $\geq t$.
- Result: most elements in \mathcal{W}_k are contained a set of size $t + 2k$.

Need to find a good condition to insert there.